

## Section 3.2 C++14

## auto Return

```

int    local2b = func2();    // Valid? Call one of the definitions of func2.
char  local2c = func3();    // OK, call to inline function func3
char  local2d = func4("a"); // OK, call to instantiation func4<const char>

// file3.cpp:
#include <file1.h>          // Error, IFNDR, redefinition of func2
auto func1() { return 1.2; } // OK, defined to return double
double local3a = func1();   // OK
int    local3b = func2();   // Valid? Call one of the definitions of func2.
char  local3c = func3();   // OK, call to inline function func3
char  local3d = func4("b"); // OK, call to instantiation func4<const char>

```

Because `func1` is declared in `file1.h` but defined in `file3.cpp`, `file2.cpp` does not have enough information to deduce its return type. Conversely, `func2` has the reverse problem: there is an ODR violation because `func2` is redefined in every translation unit that has `#include <file1.h>`. The compiler is not required to diagnose most ODR violations, but linkers will typically complain about multiply-defined public symbols. Finally, `func3` is **inline** and `func4` is a template; like `func1`, their definitions are visible in each translation unit, making the deduced return type available, but unlike `func2`, they do not create an ODR violation.

## Placeholders in trailing return types

If `auto` or `decltype(auto)` is used in a trailing return type, the meaning is the same as using the same placeholder as a leading return type:

```

auto f1() -> auto;
auto f2() -> decltype(auto);
auto f3() -> const auto&;

auto          f1(); // OK, compatible redeclaration of f1
decltype(auto) f2(); // OK, compatible redeclaration of f2
const auto&    f3(); // OK, compatible redeclaration of f3

```

When any trailing return type is specified, the leading return-type placeholder must be plain `auto`:

```

decltype(auto) f4() -> auto; // Error, decltype(auto) with trailing return
auto&          f5() -> int&; // Error, auto& with trailing return

```

## Deduced return types for lambda expressions

As described in Section 2.1. “Lambdas” on page 573, the return type of a **closure** call operator can be deduced automatically from its **return statement(s)**:

```

auto y1 = [](int& i)           { return i += 1; }; // Deduce int.

```